



INTEL IPP
REALISTIC RENDERING
MOBILE PLATFORM SOFTWARE
DEVELOPMENT KIT

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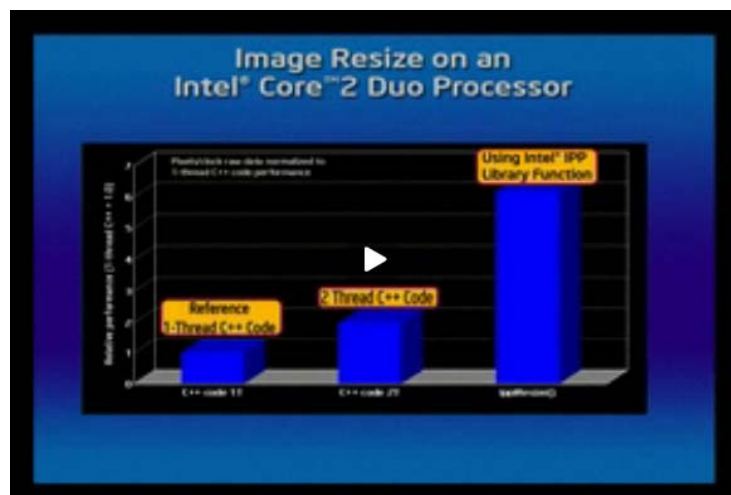
INTEL PERFORMANCE LIBRARY: INTEGRATED PERFORMANCE PRIMITIVES

- What is IPP?
 - provide optimizing software building blocks to complement Intel's optimizing compilers and performance optimization tools
- Application domain:
 - Digital media, web/enterprise data, **embedded**, communication, scientific/technical
- Target platform:
 - 32-bit Intel® architecture Platforms, 64-bit Intel 64 architecture-based platforms, **Intel XScale microarchitecture-based platforms**
- How to get IPP? – buy ☹



EMBEDDED APPLICATIONS ON INTEL XSCALE MICROARCHITECTURE-BASED PLATFORMS

	OS Version	Supported Compilers
Microsoft Windows	Windows CE 6.0	Microsoft eMbedded Visual C++* 4.0 with Service Pack 4
Linux	MontaVista Linux* 4.0 CEE LE	iwmmxt_le-gcc (MontaVista) for LE applications
	MontaVista Linux 3.1 Pro BE	xscale_be-gcc (MontaVista) for BE applications



Test code: image resizing
IPP performance in multi-core environment



PERFORMANCE-OPTIMIZED FUNCTION DOMAIN OF IPP

- Video Coding, JPEG Coding, Audio Coding, Image Processing, Speech Coding, Speech Recognition, Computer Vision, Signal Processing, Vector/Matrix Operations, Color Conversion, Data Compression, Cryptography, String Processing, Ray Tracing/Rendering (New!)
- Intel IPP functions are designed to deliver performance beyond what optimized compilers alone can deliver, by matching the function algorithms to **low-level optimizations based on the processor's available features such as Streaming SIMD Extensions (SSE, SSE2, SSE3, SSSE3, and SSE4) and other optimized instruction sets.**



KD-TREE GENERATION SAMPLE CODE SUITABLE FOR MULTI-THREAD

- Review general structure
 - Minimized kd-tree node structure
 - Multi-pass kd-tree generation for multi-thread calculation
 - IPP library function interface for kd-tree generation



```

134 /* Parallel kd tree construction using vertical subdivision */
135 IppStatus kdTreeBuild_th(IpprKdTreeNode **pDstKdTree, const Ipp32f * const pSrcVert, const Ipp32s * const pSrcTriInx,
136     Ipp32s SrcVertSize, Ipp32s SrcTriSize, Ipp32s *pDstKdTreeSize, const void * const pBldContext, int threadCnt, bool *aType)
137 {
138     const int DepthGain      = 2;
139     const int SizeThreshold  = 1000;
140     const int TrngThreshold  = 1;
141
142     if((threadCnt==1)||((SrcTriSize<SizeThreshold)||(((IpprSmp1BldContext*)pBldContext)->Alg!=ippKdBuildPureSAH)) {
143         *aType = false;
144         return ipprKdTreeBuildAlloc(pDstKdTree, pSrcVert, pSrcTriInx, SrcVertSize, SrcTriSize, pDstKdTreeSize, pBldContext);
145     }
146
147     Ipp32s          iTreeDepth = 0, rTreeDepth = 1, iTreeSize;
148     IpprPSAHBldContext iBC;
149     IpprKdTreeNode *iTreeRoot, *rTreeRoot;
150     std::vector<localSubTree> subTrees;
151     IppStatus          ipp_status = ippStsNoErr;
152
153     ippsCopy_8u((Ipp8u*)(pBldContext), (Ipp8u*)&iBC, sizeof(IpprPSAHBldContext));
154     while(rTreeDepth < threadCnt) {iTreeDepth++; rTreeDepth<<=1;};
155     iTreeDepth += DepthGain; rTreeDepth = iBC.MaxDepth - iTreeDepth;          // iTreeDepth: initial kd-tree depth, rTreeDepth: left kd-tree depth
156                                         // so, final kd-tree depth will be iBC.MaxDepth
157
158     // initial subtree with SAH kd-tree
159     /* Build initial subtree */ {
160         iBC.MaxDepth = iTreeDepth;
161         ipp_status = ipprKdTreeBuildAlloc(&iTreeRoot, pSrcVert, pSrcTriInx, SrcVertSize, SrcTriSize, &iTreeSize, &iBC);
162         if(ipp_status!=ippStsNoErr) return ipp_status;
163     }
164
165     /* Populate task queue */ {
166         Ipp32s          stack_pos = 1, dim;          // local variables
167         localSubTree   stack[MAX_KDTREE_DEPTH + 1]; // allocate stack array
168         IpprKdTreeNode *current, *left;
169
170         stack[0].node = iTreeRoot; stack[0].size = 0;          // set first stack by root node
171         ippsCopy_8u((Ipp8u*)(iBC.Bounds), (Ipp8u*)&(stack[0].bbox), sizeof(IppBox3D_32f)); // copy root bounding box
172
173         // while traversal simple kd-tree structure, find nodes which have large number of trinagles and insert them to subtree list
174         while(stack_pos!=0) {
175             stack_pos--;
176             current = stack[stack_pos].node;
177             if(current->flag_k_ofs >= 0) {          // if has child nodes
178                 left = (IpprKdTreeNode *)((long)current) + (current->flag_k_ofs)&(~0x3)); // left node offset
179                 dim = (current->flag_k_ofs)&(0x3); // split axis (last 2 bits)
180                 ippsCopy_8u((Ipp8u*)&(stack[stack_pos]),(Ipp8u*)&(stack[stack_pos+1]),sizeof(localSubTree)); // copy local sub tree stack info.
181                 stack[stack_pos].bbox[0][dim] = stack[stack_pos+1].bbox[1][dim] = current->tree_data.split; // right tree max, left tree min
182                 stack[stack_pos+1].node = left + 1; // right tree pointer
183                 stack[stack_pos+1].node = left; // left tree pointer
184             } else {
185                 if(current->tree_data.items > TrngThreshold) subTrees.push_back(stack[stack_pos]); // more then triangle threshold
186             }
187         }
188     }

```

```

134 /* Parallel kd tree construction using vertical subdivision */
135 IppStatus kdTreeBuild_th(IpprKdTreeNode **pDstKdTree, const Ipp32f * const pSrcVert, const Ipp32s * const pSrcTriInx,
136 {
137     typedef struct KdTreeNode{
138         Ipp32s  flag_k_ofs;
139         union _tree_data{
140             Ipp32f  split;
141             Ipp32s  items;
142         }tree_data;
143     }IpprKdTreeNode;
144     return ipprKdTreeBuildAlloc(pDstKdTree, pSrcVert, pSrcTriInx, SrcVertSize, SrcTriSize, pDstKdTreeSize, pBldContext);
145 }
146 typedef struct SimpleBuilderContext{
147     IpprKdTreeBuildAlg  Alg;
148     Ipp32s              MaxDepth;
149 }IpprSmp1BldContext;
150 where Alg - must be equal to ipprKdTreeBuildSimple constant;
151 Ipp32s MaxDepth - reserved.
152
153 typedef enum {
154     wh ipprKdTreeBuildSimple = 0x499d3dc2, // Simple building mode
155     it ipprKdTreeBuildPureSAH = 0x2d07705b // SAH building mode
156 }IpprKdTreeBuildAlg;
157
158 // typedef struct PSAHBuilderContext{
159 /*
160     IpprKdTreeBuildAlg  Alg;
161     Ipp32s              MaxDepth;
162     Ipp32f              QoS;
163     Ipp32s              AvailMemory;
164     IppBox3D_32f       *Bounds;
165 */
166 }IpprPSAHBldContext;
167 where Alg - must be equal to ipprKdTreeBuildPureSAH constant;
168     MaxDepth - maximum tree subdivision depth (minimum - 0, maximum - 50);
169     QoS - termination criteria modifier (minimum - 0.0, maximum - 1.0);
170     AvailMemory - maximum available memory in Mb;
171     *Bounds - cut-off bounding box.
172
173 // while traversal simple kd-tree structure, find nodes which have large number of trinagles and insert them to subtree list
174 while(stack_pos!=0) {
175     stack_pos--;
176     current = stack[stack_pos].node;
177     if(current->flag_k_ofs >= 0) {
178         left = (IpprKdTreeNode *)((long)(current) + (current->flag_k_ofs)&(~0x3)); // if has child nodes // left node offset
179         dim = (current->flag_k_ofs)&(0x3); // split axis (last 2 bits)
180         ippsCopy_8u((Ipp8u*)&(stack[stack_pos]),(Ipp8u*)&(stack[stack_pos+1]),sizeof(localSubTree)); // copy local sub tree stack info.
181         stack[stack_pos].bbox[0][dim] = stack[stack_pos+1].bbox[1][dim] = current->tree_data.split; // right tree max, left tree min
182         stack[stack_pos+1].node = left + 1; // right tree pointer
183         stack[stack_pos+1].node = left; // left tree pointer
184     } else {
185         if(current->tree_data.items > TrngThreshold) subTrees.push_back(stack[stack_pos]); // more then triangle threshold
186     }
187 }
188 }

```

tial kd-tree depth, rTreeDepth: left kd-tree depth
ree depth will be iBC.MaxDepth

cTriSize, &iTreeSize, &iBC);

// set first stack by root node
); // copy root bounding box

```

190  /* Threaded build of subtrees (75% of total time) */ {
191      Ipp32s          i = 0, leaf_cnt = subTrees.size();
192      IppStatus      ipp_status = ippStsNoErr;
193  #if defined (_OPENMP)
194  #pragma omp parallel num_threads(threadCnt) default(shared)
195  #endif
196      {
197          Ipp32s          cIdx;
198          IpprPSAHBldContext lBC;
199          localSubTree    *subTree;
200          IppStatus      l_ipp_status;
201
202          ippCopy_8u((Ipp8u*)(pBldContext), (Ipp8u*)&lBC, sizeof(IpprPSAHBldContext));
203          lBC.MaxDepth = rTreeDepth;
204
205          while((i < leaf_cnt)&&(ipp_status==ippStsNoErr)) {
206  #if defined (_OPENMP)
207              #pragma omp critical
208  #endif
209              { cIdx = i++; }
210
211              if(cIdx < leaf_cnt) {
212                  subTree = &(subTrees[cIdx]);
213                  subTree->node = NULL; lBC.Bounds = &(subTree->bbox);
214                  l_ipp_status = ippKdTreeBuildAlloc(&(subTree->node), pSrcVert, pSrcTriInx, SrcVertSize, SrcTriSize, &(subTree->size), &lBC);
215                  if(l_ipp_status!=ippStsNoErr) ipp_status = l_ipp_status;
216              }
217          }
218      }
219  }

```



Due to the algorithm specific implementation, initial memory allocation for SAH-based tree building can not be less than 80Mb even for very small scenes, this limits a minimal useful value of the AvailMemory by 81Mb.


```

221  /* Subtrees merge cycle (0.05% of total time) */ {
222  if(ipp_status==ippStsNoErr) {
223      IpprKDTreeNode *n1, *n2, *n3, *n4;
224      Ipp32s *triPtr, node_cnt = 2, trng_cnt = 0, st_index = 0;
225
226      // count subtree nodes and leaf node triangles|
227      n1 = iTreeRoot + 2; node_cnt+=localFindLastNode(n1, &n2); // find last(right end) node of the tree
228      for(;n1<n2;n1++) { // search whole tree
229          if((n1->flag_k_ofs < 0)&&(n1->tree_data.items > 0)) { // if leaf node
230              if(n1->tree_data.items > TrngThreshold) { // if has large number of triangles
231                  n3 = subTrees[st_index++].node + 2; node_cnt += localFindLastNode(n3, &n4); // then there are subtree nodes
232                  for(;n3<n4;n3++) if((n3->flag_k_ofs < 0)&&(n3->tree_data.items > 0)) trng_cnt += n3->tree_data.items;
233              } else { trng_cnt += n1->tree_data.items; }
234          }
235      }
236
237      st_index = node_cnt*sizeof(IpprKDTreeNode) + trng_cnt*sizeof(int);
238      rTreeRoot = (IpprKDTreeNode *)ippMalloc_02s(st_index>>2);
239      if(rTreeRoot==NULL) {
240          ipp_status = ippStsNoMemErr;
241      } else {
242          IpprKDTreeNode *hBTnTarget = rTreeRoot + 2;
243          Ipp32s *hBTiTarget = (Ipp32s *) (rTreeRoot + node_cnt),
244              hSubTreeId = 0;
245
246          rTreeRoot[1].flag_k_ofs = 0x00000000;
247          rTreeRoot[1].tree_data.items = 0x00000000;
248
249          localTreePacking(iTreeRoot, subTrees, TrngThreshold, rTreeRoot, &hBTiTarget, &hBTnTarget, &hSubTreeId);
250
251          *pDstKDTree = rTreeRoot;
252          *pDstKDTreeSize = st_index;
253          *aType = true;
254      }
255  }
256 }
257
258 /* Free intermediate subtrees */ {
259 for(int i=0; i<subTrees.size(); i++) {
260     if(subTrees[i].size!=0) ipprKDTreeFree(subTrees[i].node);
261 }
262 ipprKDTreeFree(iTreeRoot);
263 }
264
265 return ipp_status;
266 }

```

```

120 long localFindLastNode(IpprKDTreeNode* root, IpprKDTreeNode** last){
121     IpprKDTreeNode *n1, *n2;
122
123     n1 = root; n2 = n1 + 1;
124     while((n1->flag_k_ofs >= 0) || (n2->flag_k_ofs >= 0)) {
125         if(n1->flag_k_ofs >= 0) n1 = (IpprKDTreeNode *)((long)(n1) + (n1->flag_k_ofs)&(~0x3));
126         if(n2->flag_k_ofs >= 0) n2 = (IpprKDTreeNode *)((long)(n2) + (n2->flag_k_ofs)&(~0x3));
127         if(n2 < n1) { n2 = n1 + 1; } else { n1 = n2; n2++; }
128     }
129     *last = n2 + 1;
130
131     return (*last - root);
132 }

```

```

83 /* Merge subTrees into top level pRoot tree instead of leafs which contain more the Threshold triangles */
84 void localTreePacking(const IpprKDTreeNode *pRoot, std::vector<localSubTree> &subTrees, Ipp32s Threshold, IpprKDTreeNode *hBTRoot,
85 Ipp32s **hBTiTarget, IpprKDTreeNode **hBTnTarget, Ipp32s *hSubTreeId)
86 {
87     if(pRoot->flag_k_ofs >= 0) { // if has child node
88         IpprKDTreeNode *node, *leftSon;
89         Ipp32s dim;
90
91         node = (IpprKDTreeNode *)((long)(pRoot) + (pRoot->flag_k_ofs)&(~0x3)); // get child node
92         dim = (pRoot->flag_k_ofs)&(~0x3); // get split dimension
93         hBTRoot->tree_data.split = pRoot->tree_data.split; // set split value
94         hBTRoot->flag_k_ofs = ((long)(*hBTnTarget) - (long)(hBTRoot)) | dim; // set node offset
95         leftSon = *hBTnTarget; *hBTnTarget += 2;
96
97         localTreePacking(node, subTrees, Threshold, leftSon, hBTiTarget, hBTnTarget, hSubTreeId);
98         localTreePacking(node + 1, subTrees, Threshold, leftSon + 1, hBTiTarget, hBTnTarget, hSubTreeId);
99     } else { // if leaf node
100         Ipp32s size = pRoot->tree_data.items;
101
102         if(size < 0) {
103             hBTRoot->flag_k_ofs = pRoot->flag_k_ofs;
104             hBTRoot->tree_data.items = pRoot->tree_data.items;
105         } else {
106             if((size > Threshold)&&(Threshold >= 0)) { // if has large number of tris
107                 Ipp32s stNum = (*hSubTreeId)++; // then there will be subtree
108
109                 localTreePacking(subTrees[stNum].node, subTrees, -1, hBTRoot, hBTiTarget, hBTnTarget, hSubTreeId);
110             } else {
111                 Ipp32s *pTriIdx = (Ipp32s*)((long)pRoot - pRoot->flag_k_ofs); // node offset - positive integer
112                 // triangle offset - negative integer
113                 hBTRoot->tree_data.items = size; // set number of triangles
114                 hBTRoot->flag_k_ofs = (long)(hBTRoot) - (long)(*hBTiTarget); // set triangle index offset
115                 while(size--) *(*hBTiTarget)++ = *(pTriIdx++); // copy triangle index
116             }
117         }
118     }
119 }

```

INTEL MOBILE PLATFORM SOFTWARE DEVELOPMENT KIT

- An Open Source Project for Windows* and Linux*, and Moblin(mobile linux internet)
- Quickly produce efficient, reliable, cross platform context-aware mobile applications that are tuned for the mobile computing environment.
- Close the gap between software functionality and mobile device platforms



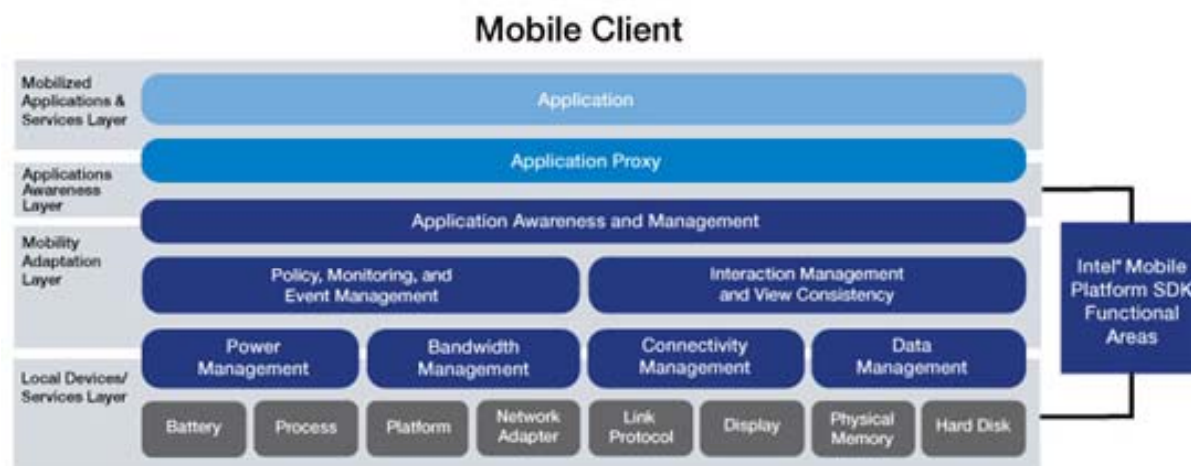
FEATURES OF IMP-SDK

- **Manage connectivity transparently.** Users can focus on their tasks rather than managing network connections.
- **Effectively balance power and performance.** Users can get aware of and utilize the available battery life.
- **Work across multiple platforms.** Users can access applications on the device of their choice, taking advantage of device capabilities.
- **Use available memory and disk space for local data store and synchronization.** Applications developed with the Mobile Platform SDK 1.3 can discover and utilize available memory and disk storage for caching, local data store, and synchronization.
- **Adapt to different display types.** Mobile devices are often connected to different display types. Applications based on the Mobile Platform SDK 1.3 can discover the attached display type and adapt to it.
- **Manage network bandwidth.** Control the bandwidth of network traffic in system, application, process and socket level.



THE MOBILITY SOFTWARE CHALLENGE

- Creating applications that are aware of platform context and resources in order to efficiently take advantage of mobile platforms
- Developing cross-platform, cross-runtime solutions for application deployment across multiple clients



Virus Scanner Example

